# Uniform Random Sampling of Chemical Space

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#### Coverage

- Systematic databases are small
- Imbalanced
- Synthesis creates biases



ChEMBL	
[0,5)	250
[5,10)	5.939
[10,15)	50.802
[15,20)	231.873
[20,25)	535.187
[25,30)	605.591
[30,35)	457.611
[35,40)	236.646
[40,45)	106.679
[45,50)	48.321
[50,55)	27.320
[55,60)	16.730
[60,65)	13.014
[65,70)	8.599
[70,75)	2.327
[75,79]	16

## Random Sampling

Allow for data-driven fundamental statements "Most molecules do X", "High X means low Y"

**Transferability** More reliable understanding of trends

Lower data bias More realistic generalisation error

More data efficiency Maximally spanning coverage

**Formal statements** Often require uniform sampling

Measure coverage Generative methods



#### Problems

- Total number unknown
- Distribution unknown

Scaling

#### Take some atoms from {C, O, N, F, S}, H-saturated







"Give me some fruit, please!"

## Random Sampling

#### Goal

Sample all molecules (with given constraints) with known probabilities.

#### Sampling

- Choose weighted random chemical formula
- Choose weighted random degree sequence
- Choose weighted random molecular graph

#### Requirements

- Find all sum formulas and degree sequences
- Sample loop-free multigraphs with given degree sequences uniformly
- Find weights





Solved, Seconds Solved, Seconds<sup>[1]</sup>

1 C Greenhill, M Sfragara, *Theo Comput Sci*, **2018**.

#### Counting via enumeration

- SMOG (1996), MOLGEN (1998), ASSEMBLE(2000), OMG (2012), PMG (2013), MAYGEN (2021), surge (2022)
- Until about 10-15 atoms

#### Orderly generation

- Find canonical sorting of (partial) molecular graphs
- Create graphs in canonical order

#### Monte Carlo Sampling

- Grow and shrink molecular graphs
- Slow sampling

# Counting without enumeration

#### Goal

Estimate number n of loop-free multigraphs with given degree sequence.

#### Average Path Length l<sub>G</sub>

Sample from random molecule pairs

 $l_G \sim \log n$ 





Pure degree sequenceEvery valency exists only once.O=C=SO=C=OS=C=S

#### Non-pure estimate

- Assumes that random graphs are almost never symmetric
- All modifications independent
- Combinatorial product

$$N_P(d) = \prod_v \prod_i \left( \sum_{j>i} c_j \atop c_i \right)$$
$$l_G(d) = \left( 1 + \left[ \sum_i d_i \right]^{-1} \log N_P^L \right) l_G(d_U)$$



### Asymptotic Scaling

#### Average Path Lengths become expensive

- Heuristics become less efficient
- More sampling
- Converging slowly

#### Asymptotic Scaling<sup>[1]</sup>

- Needs to be calibrated to molecules



 $G = \frac{M!}{(M/2)!2^{M/2}k_1!\cdots k_n!}$  $\exp\left(\left(y_1 - \frac{1}{2}\right)\frac{M_2}{M} + \left(x_2 - \frac{1}{2}\right)\frac{M_2^2}{2M^2} + \frac{M_2^4}{4M^5}\right)$  $-\frac{M_2^2 M_3}{2M^4} + \left(x_3 - x_2 + \frac{1}{3}\right) \frac{M_3^2}{2M^3} + (an+b)/M + (cn+d)M + e$  $M_r = \sum_{i}^{n} [k_i]_r$ 

1 C Greenhill, B McKay, *SIAM*, **2013**.

### Examples

#### 10 atoms, CHONF, at least 3 hydrogens and one fluorine

CH<sub>2</sub>

F

11



- Number of atoms
- Atoms per element (and combinations)
- Valences
- Via rejection sampling:
  - Bonds (count and bond orders)
  - Ring presence / membership
  - Stability
  - Substructures

- ...



### Coverage of databases



#### Chemical formula

### Summary



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Randomized | Known distribution: statistical statements
Regions | No one-by-one iteration
Bias reduction | Datasets and predictive power
Seeding | Maximally spanning datasets or Monte-Carlo acceleration



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